Comparing Quotient- and Symmetric Containers

Philipp Joram Niccolò Veltri

Tallinn University of Technology, Estonia

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Containers: syntax for polymorphic programs

We want to study various notions of containers in HoTT.

- containers define syntax for (polymorphic) data types.
- morphisms of containers are syntax for transformations of such data
- ▶ A functor $\llbracket \rrbracket$ gives semantics to containers & their morphisms
- \blacktriangleright [[-]] often lands in endofunctors and natural transformations

Today:

Quotient containers, and how they relate to symmetric containers.

Quotient containers

Idea: Data at certain positions are identified. A quotient container $(S \triangleleft P/G)$ is...

- ▶ a set of *shapes S*,
- a set of *positions* P(s) for each shape s
- ▶ a group of permissible symmetries $G_s \leq \operatorname{Aut}(P(s))$

Extension of a quotient container: an endofunctor $[S \triangleleft P/G]_{/}$: Set \rightarrow Set:

$$\llbracket S \triangleleft P/G \rrbracket/(X) := \sum_{s:S} \frac{P(s) \to X}{\sim_s} \qquad u \sim_s v :\iff \frac{P(s) \xrightarrow{\exists g:G_s}}{\bigvee} P(s)$$

Example: unordered pairs

The container of *unorderd* pairs has

$$JPair := (1 \triangleleft 2 / Aut(2))$$

$$two positions \uparrow \dots that identify data by swapping$$

The extension of UPair gives sets of unordered pairs:

$$\llbracket \mathsf{UPair} \rrbracket_{/}(X) = \sum_{s:1} (\mathbf{2} \to X) / \sim_s = \frac{X^2}{(x,y) \sim (y,x)}$$

More examples: Finite multisets, cyclic lists. Non-examples: Finite sets.

Category of quotient containers

Quotient containers¹ are the objects of a category \mathcal{Q} . Containers are related by *premorphisms*:

$$(\begin{array}{c} u \\ s \\ \hline \\ g(us) \\ \hline \\ elated positions stay related \end{array}) : (S < P/G) \rightarrow (T < Q/H)$$

Morphisms of \mathcal{Q} are premorphisms quotiented by some relation.

¹Abbott et al., "Constructing Polymorphic Programs with Quotient Types".

We give a definition of Q in HoTT:

- We assume shapes and positions to be (homotopy) sets, not arbitrary types
- Define extension functor using set-quotients (i.e. via a HIT)

Quotients of sets are nice to work with 2 in HoTT, so. . .

² Terms and conditions apply

Properties of quotient containers

Some immediate results, for example:

- ▶ Q-Iso(UPair, UPair) is contractible
- (UPair = UPair) $\simeq 2$
- ▶ Q is *not* a univalent category

With a little more work, we can port traditional proofs to HoTT:

Theorem

Each $\llbracket Q \rrbracket$ is a left Kan extension.

Universal property of Kan extensions implies that $[-]_{/}$ is fully faithful.

These definitions and proofs are quite involved. There is entirely too much reasoning about symmetry groups. HoTT promises to help study such symmetries. Questions:

- 1. Can we find a more intensional presentation of containers with symmetries?
- 2. Are quotient containers a subclass of those?



- ▶ morphisms are *what you expect*[™]
- define a univalent bicategory S (2-cells = homotopies of morphisms)
- $\llbracket S \triangleleft P \rrbracket(X) := \sum_{s} P(s) \rightarrow X$ is a pseudofunctor Gpd \rightarrow Gpd

³Gylterud, "Symmetric Containers".

Unordered pairs, take II

The symmetric container of *unordered pairs* is UPair := $(B(Aut(2)) \triangleleft U)$, where

▶ B(Aut(2)) has one point • and one non-trivial path (swap : • = •)

► *U* is defined by induction:

$$U(\bullet) := \mathbf{2}$$

$$U(swap) := \mathbf{ua not}$$

$$\uparrow a path \mathbf{2} = \mathbf{2}$$

The path space of shapes encodes all symmetries! For x, y : [UPair](X):

$$(x = y) = \sum_{\sigma: Aut(2)} snd(x) = snd(y) \circ \sigma$$

Delooping construction

This gives us an idea how to associate a symmetric container to any quotient container.

Definition

Define $\mathbf{B}(S \triangleleft P/G) := (S^{\uparrow} \triangleleft P^{\uparrow})$ where

$$S^{\uparrow} := \sum_{s:S} \mathbf{B}G_s$$

 $P^{\uparrow} := egin{cases} (s:S) & \mapsto P(s) : \mathcal{U} \ (g:G_s) & \mapsto \mathsf{ua}(g) : P(s) =_{\mathcal{U}} P(s) \end{cases}$

Note: A version of P^{\uparrow} appears in the definition of $[-]_{/}$ as a Kan extension.

 $[\![\mathbf{B}(Q)]\!]$ maps groupoids to groupoids. How does it relate to $[\![Q]\!]_/$? Theorem

Extension of $\mathbf{B}Q$ II

Proof.

Idea: S^{\uparrow} has pointed connected components.

$$\| \mathbb{B}Q \|(X) \| \simeq \| \sum_{s:S} \sum_{g: \mathbf{B}(G_s)} P^{\uparrow}(s,g) \to X \|$$

set truncation

$$\simeq \sum_{s:S} \| \sum_{g: \mathbf{B}(G_s)} P^{\uparrow}(s,g) \to X \| \qquad (S \text{ is a set})$$

$$\simeq \sum_{s:S} \frac{P(s) \to X}{\sim_{G_s}} \qquad (by \text{ computation})$$

$$\simeq \| \mathbb{Q} \|_{/}(X) \qquad \Box$$

Does B(-) extend to an action on Q-morphisms? Not obviously:

- ▶ maps of shapes $\sum_s \mathbf{B} G_s \to \sum_t \mathbf{B} H_t$ require a group homomorphism $G_s \to H_t$
- Q-morphisms are equivalence classes
- For a premorphism (u, f, φ) : (S ⊲ P/G) → (T ⊲ Q/H), φ : G_s → H_{ut} is not necessarily a group homomorphism

This means that B(-) cannot be directly defined by induction on Q-morphisms.

Remedies

Work harder:

- ▶ give alternative presentation of Q-morphisms
- get rid of quotiented homsets this way
- maybe it is functorial after all?

Change definitions:

- add properties: make premorphisms preserve group structure of symmetries
- all practical examples satisfy this property
- is this a parametricity condition?

More questions

Can we go the other way $\boldsymbol{\mathsf{U}}:\mathcal{S}\to\mathcal{Q}?$

- Yes, assuming a form of choice: "Every groupoid has *pointed* connected components."
- ▶ In this case, **U** is a retract of **B**: U(B(Q)) = Q.

Question to TYPES:

Is the above a known choice principle?

Conclusion

- Quotient containers are symmetric containers, in some way
- Less clear how morphisms relate
- Some interesting questions of reverse mathematics arise
- Cubical Agda has been helpful in figuring this all out

Thank you!